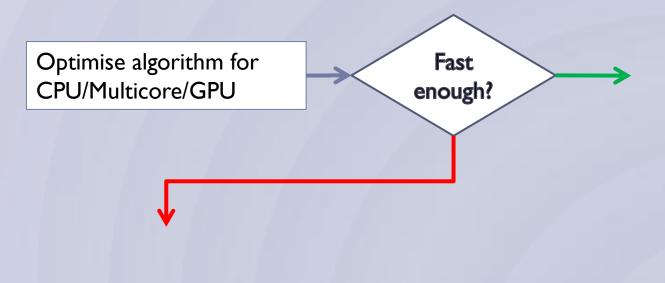
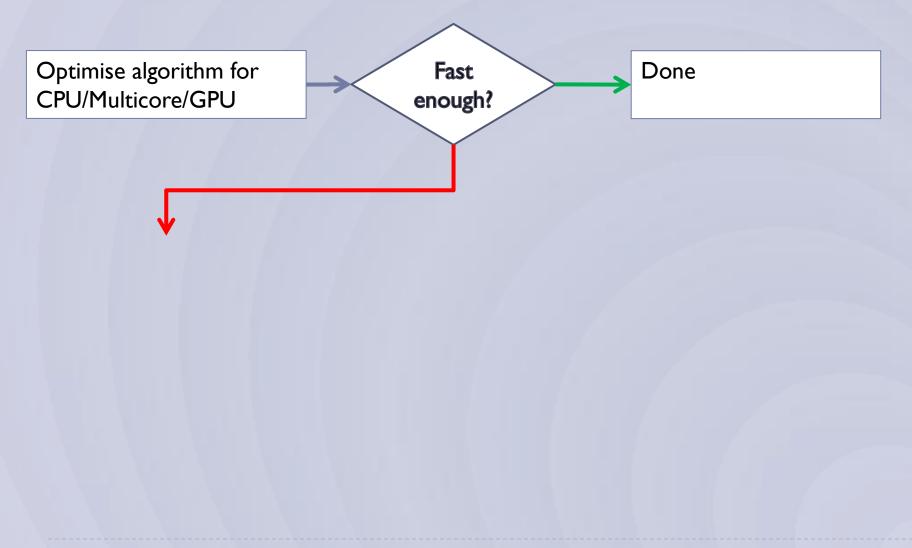
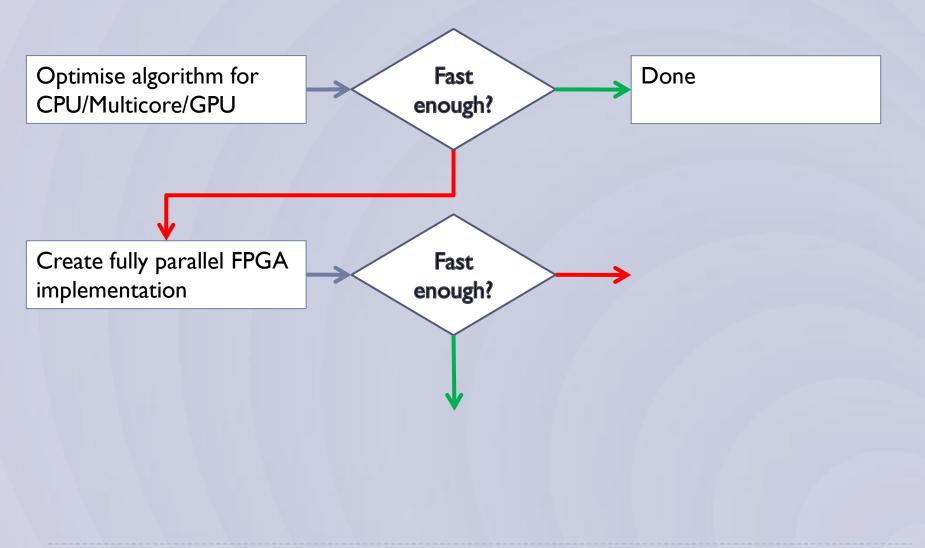
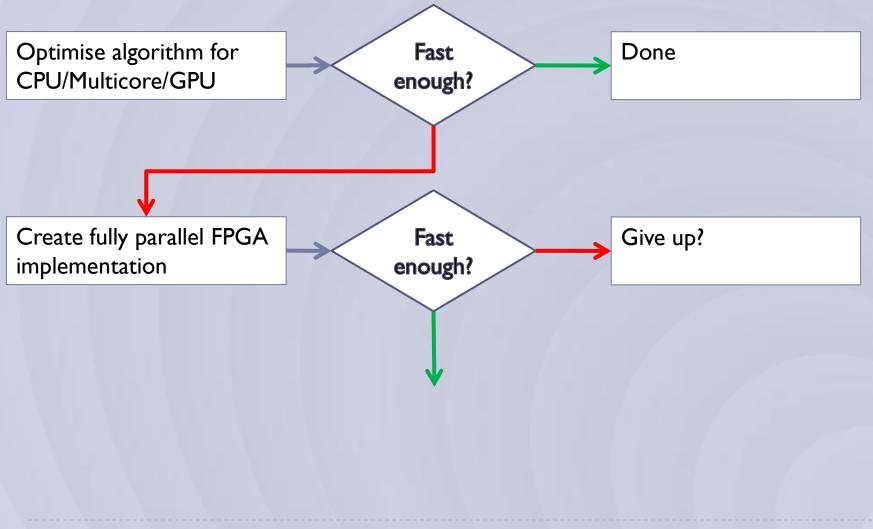
#### A Scalable Approach for Automated Precision Analysis

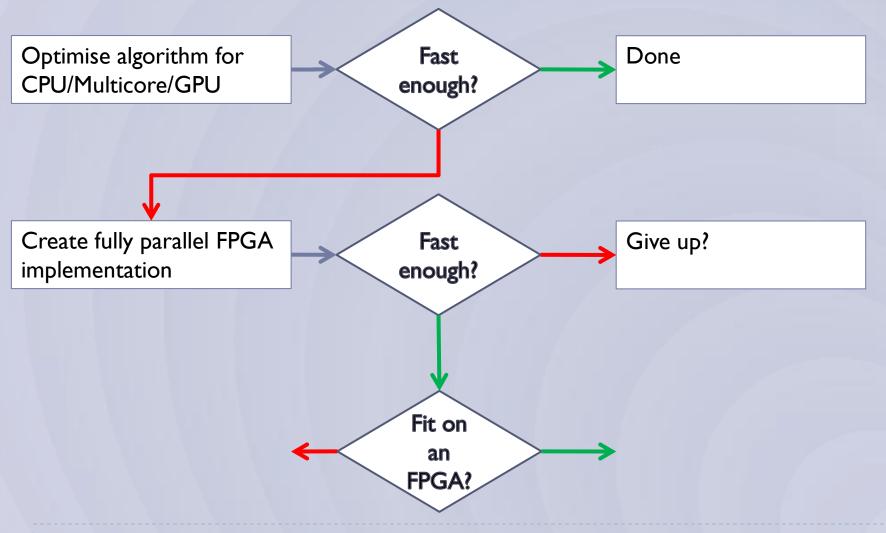
David Boland and George A. Constantinides

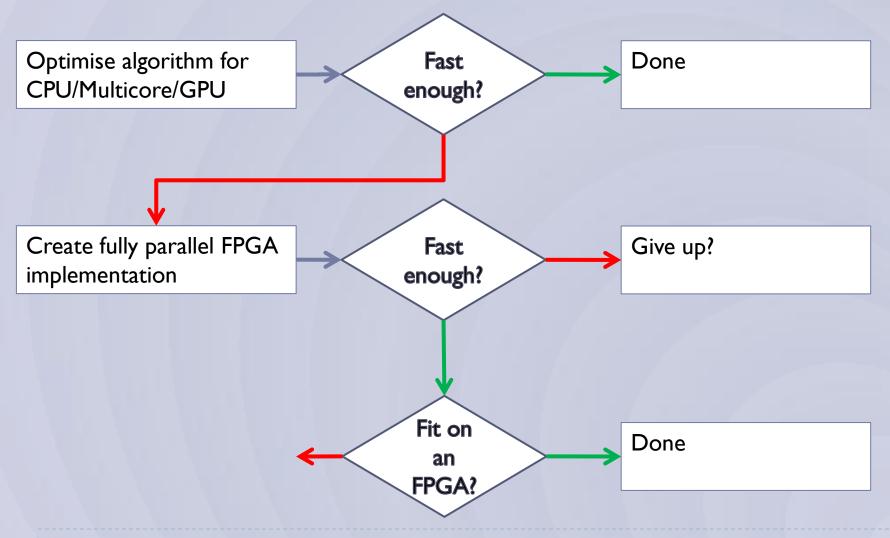


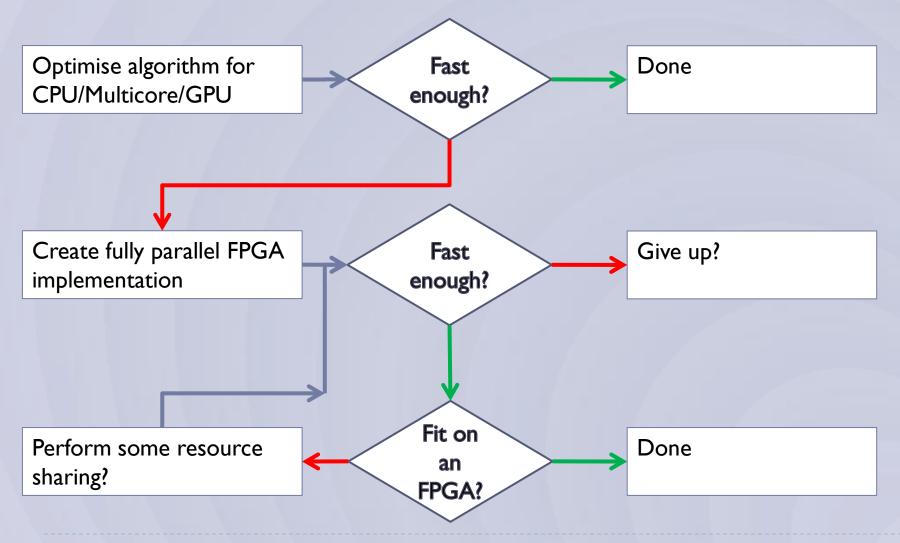


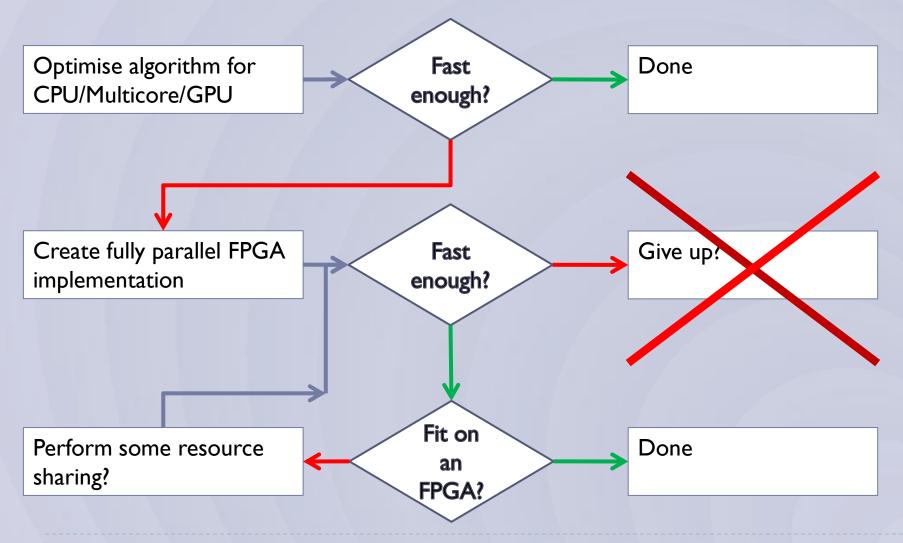












2

Word-length optimisation can be the game changer...

- Performance gain by moving from IEEE 754 double precision to single precision:
  - > 2x for a CPU
  - > 2-9x for a GPU
- FPGAs have much greater flexibility
  - Can implement any custom precision
    - Large performance trade-offs
    - Many factors affected
      - □ Silicon area
      - Clock speed
      - □ Latency
      - □ Memory use
      - Data transfer overhead

#### Reducing word-length can cause errors

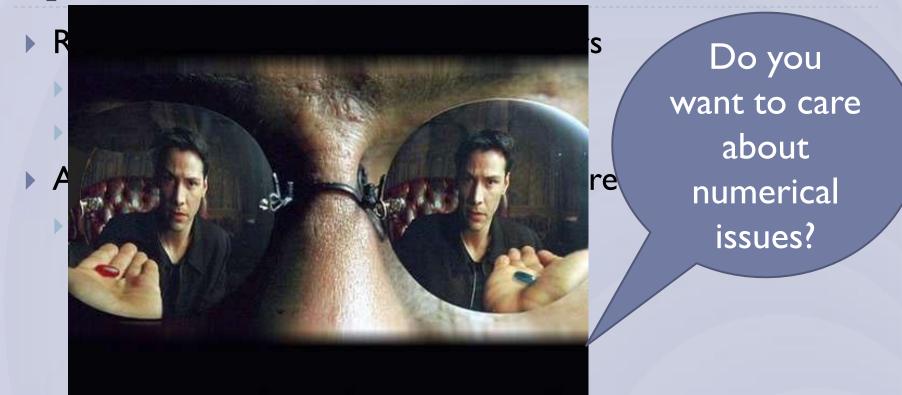
- Overflow error
- Accumulation of individual round-off errors

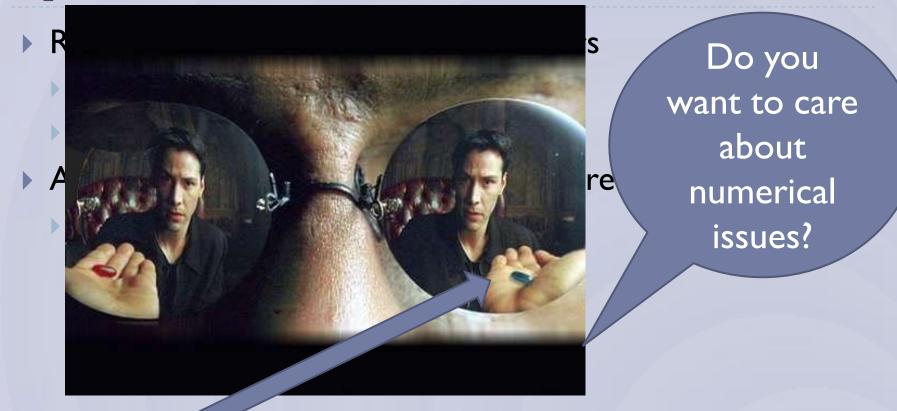
#### Reducing word-length can cause errors

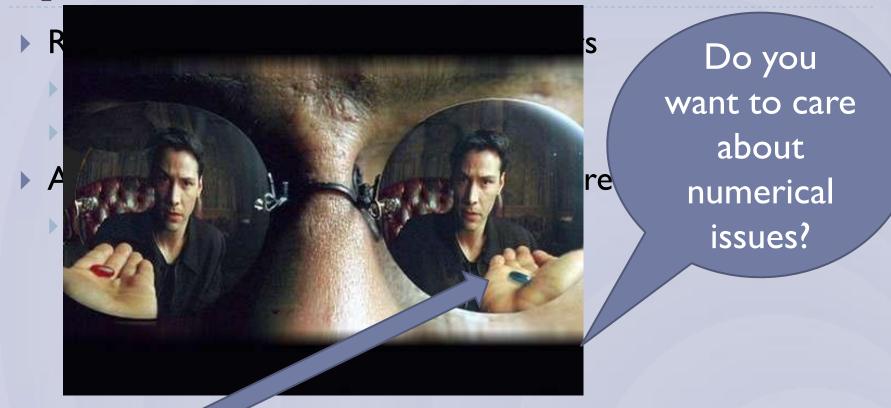
- Overflow error
- Accumulation of individual round-off errors
- Allows 'fair' comparison versus software

#### Reducing word-length can cause errors

- Overflow error
- Accumulation of individual round-off errors
- Allows 'fair' comparison versus software
  - Lazy (& incorrect?) comparison
    - (a+b)+c  $\neq$  a+(b+c)

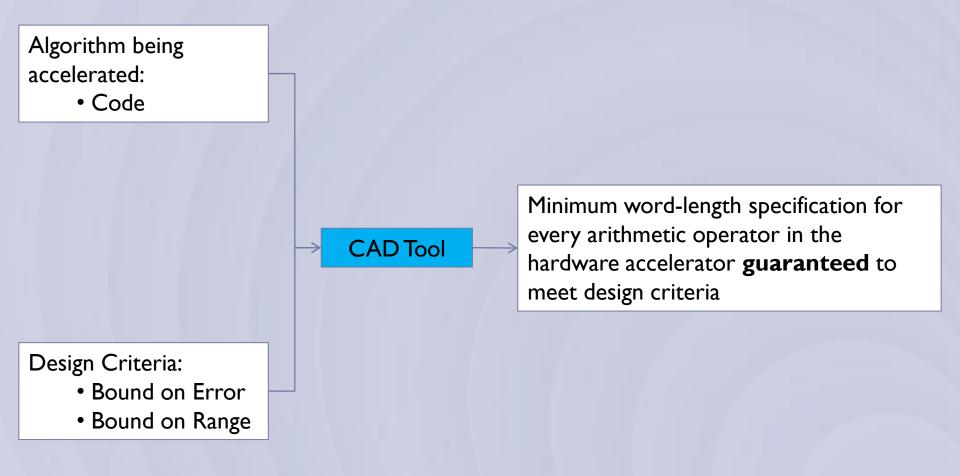


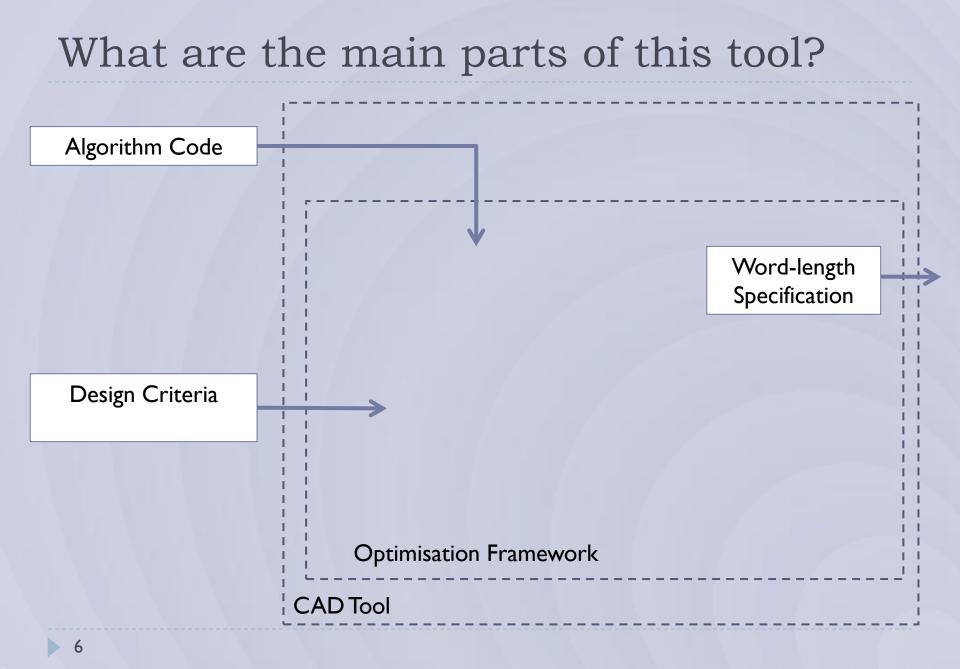


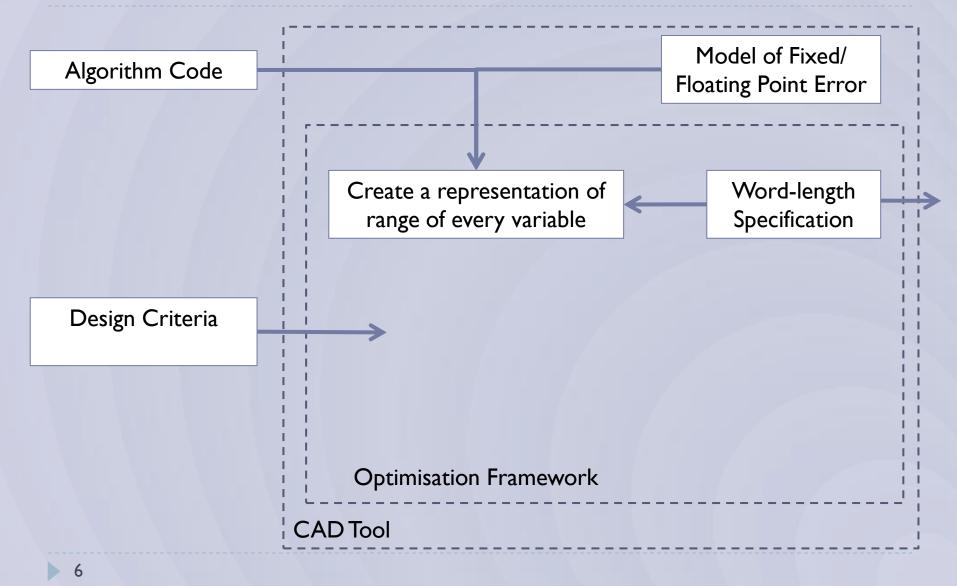


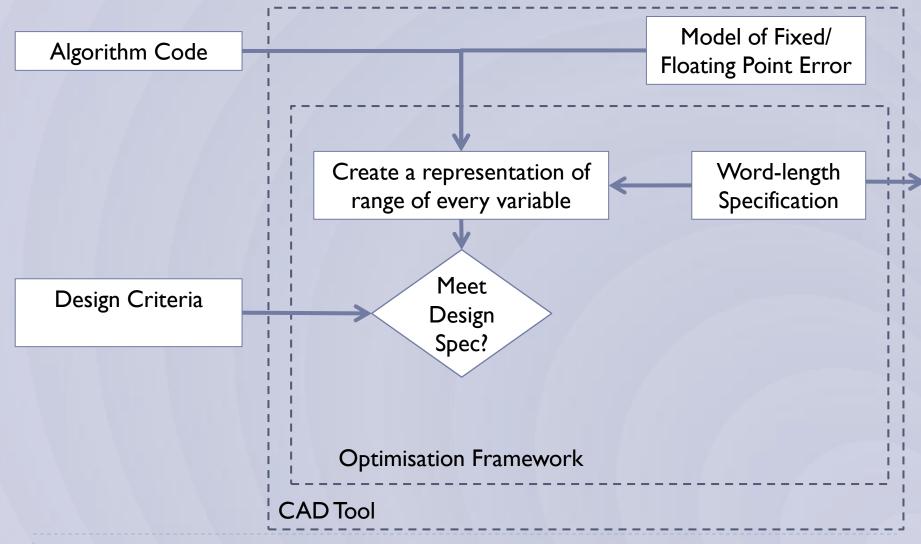
Report greater speed up factors by using IEEE single precision.

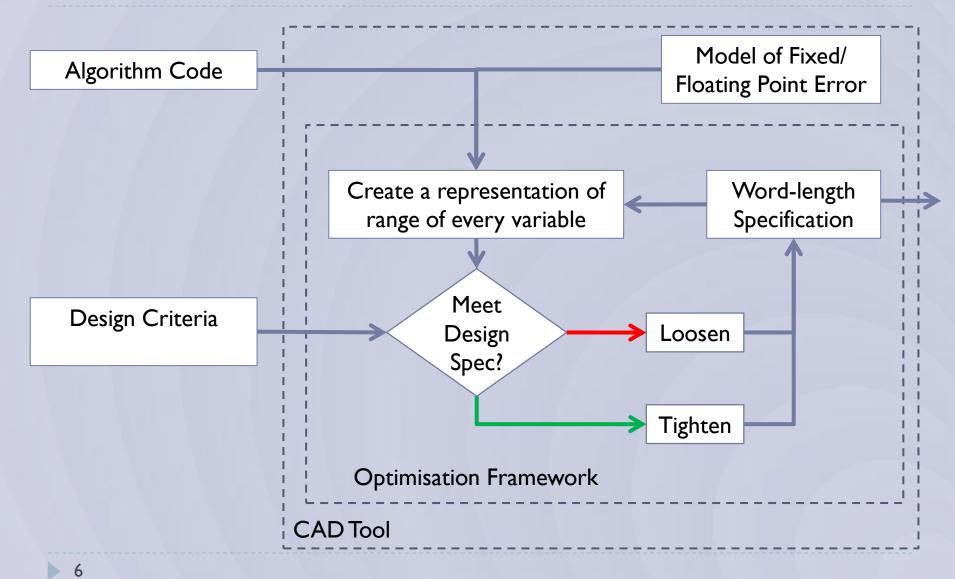
# Ideal word-length optimisation

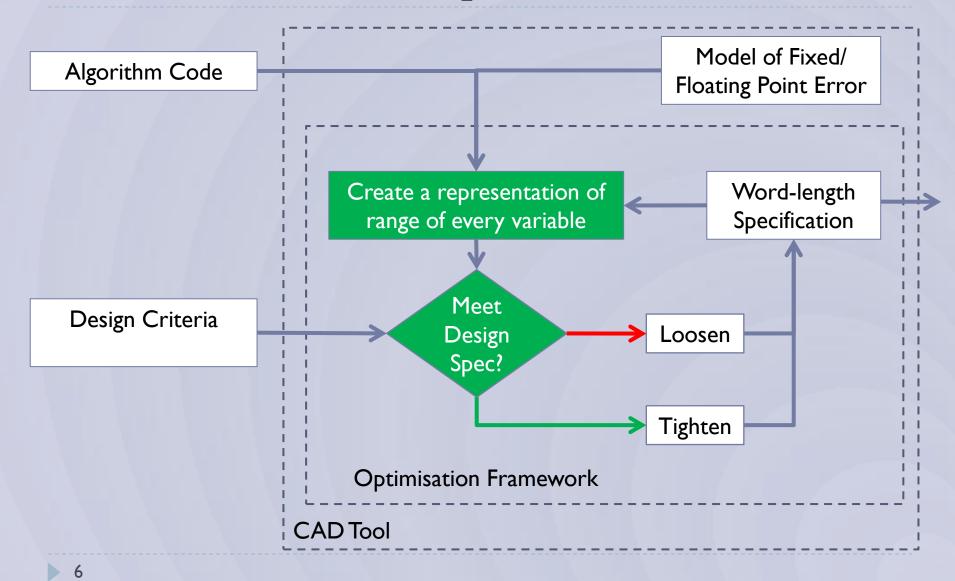














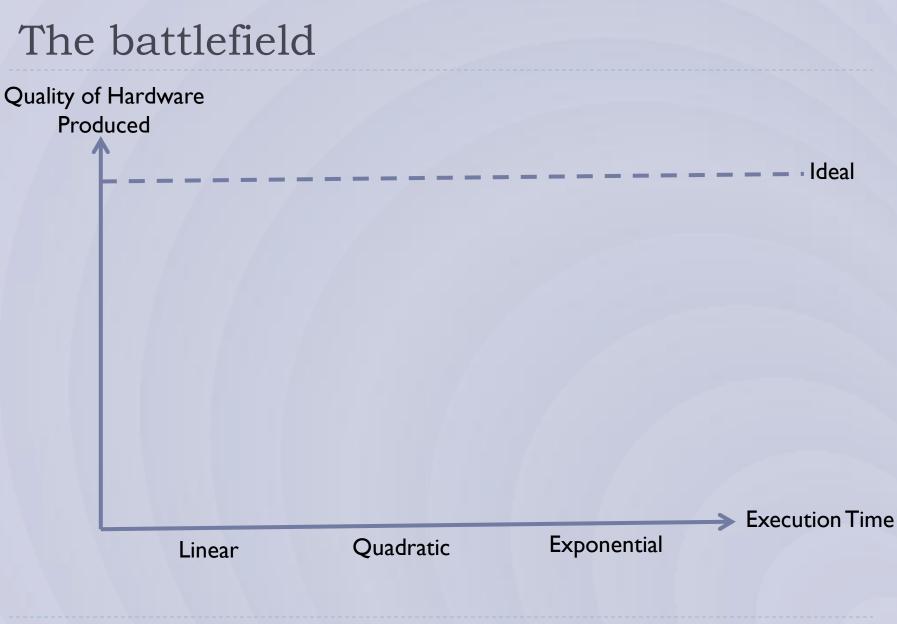
Quality of Hardware Produced

Linear

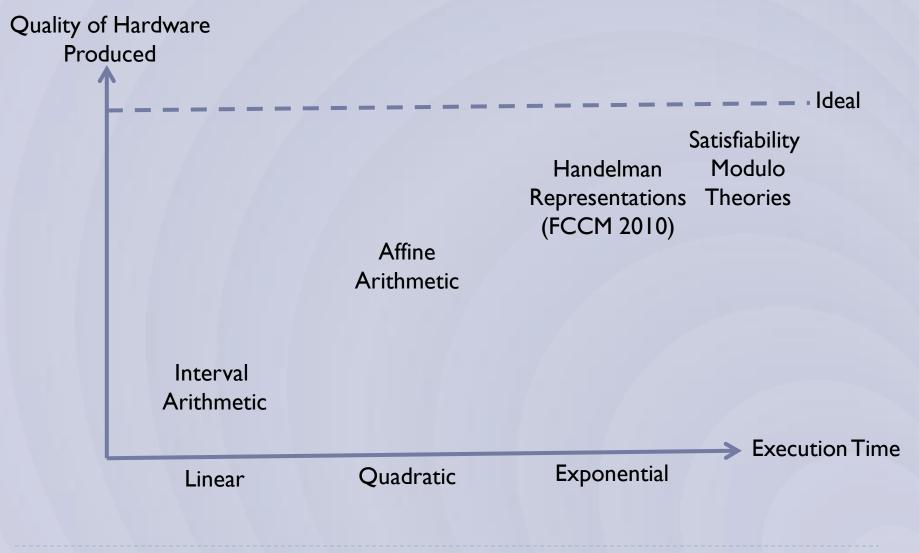
Quadratic

Exponential

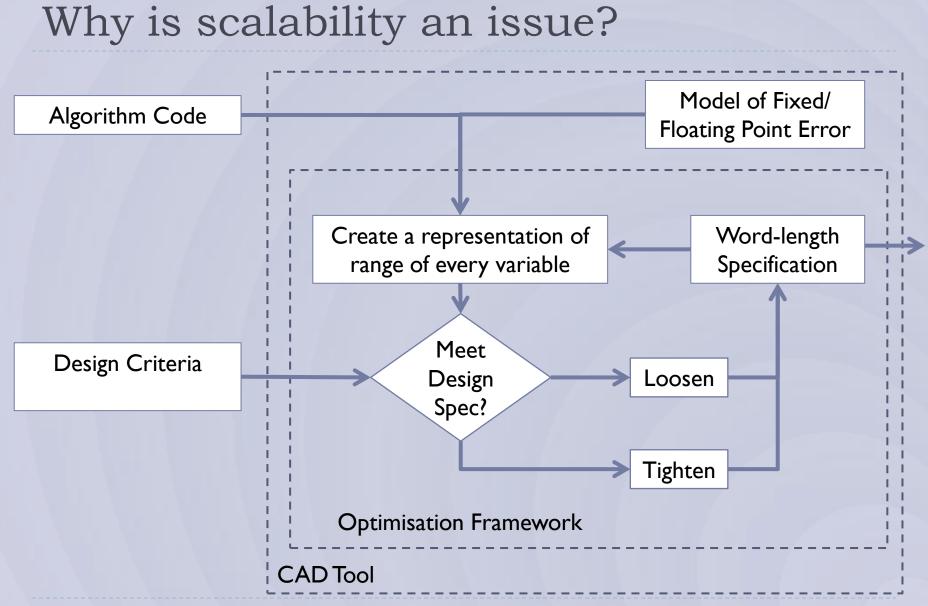
Execution Time



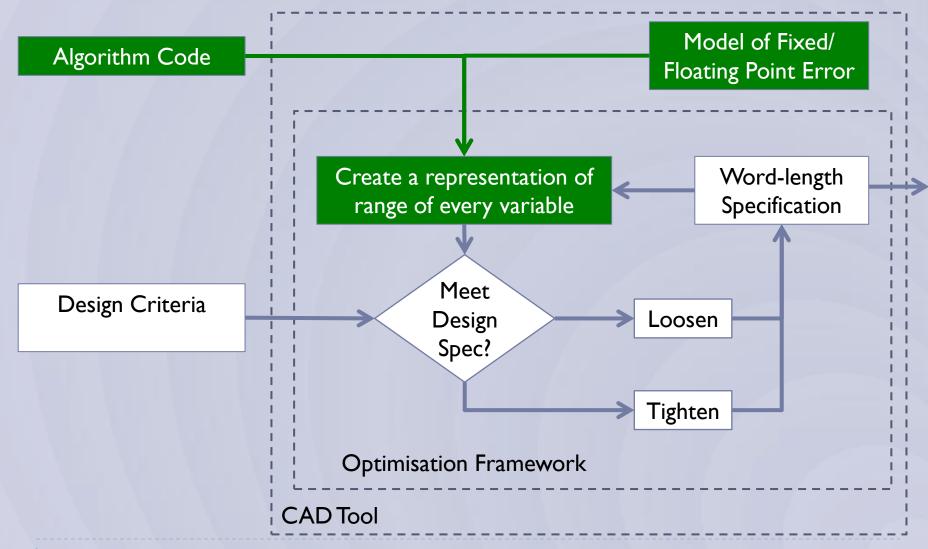
#### The battlefield



#### The battlefield Quality of Hardware Produced Ideal Satisfiability Modulo This work Handelman Theories Representations (FCCM 2010) Affine Arithmetic Interval Arithmetic Execution Time Exponential Quadratic Linear



#### 



Modelling Floating Point Error

The closest floating point approximation  $\hat{x}$  of x can be expressed as:

 $\hat{x} = x(1 + \delta_1)$   $|\delta_1| \le 2^{-m}$  (m = # of mantissa bits)

The floating point result of any scalar operation  $\bigcirc$ , where  $\bigcirc \in \{+, -, \times, \div\}$  can be bounded as:

 $\widehat{x \odot y} = (x \odot y) (1 + \delta_1)$ 

#### Simple example:

- Code: a = x \* y; b = a \* z;
- Where:  $x \in [0.8, 1.2], y \in [0.9, 1.1], z \in [9.9, 10.1]$

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- If we denote:  $|x_1| \le 0.2$ ,  $|y_1| \le 0.1$ ,  $|z_1| \le 0.1$   $|\delta_i| \le 2^{-12}$
- Then:  $x = (1 + x_1), y = (1 + y_1), z = (10 + z_1)$

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Create polynomials:

$$a = (1 + x_1)(1 + y_1)(1 + \delta_1)$$
  

$$a = 1 + x_1 + y_1 + x_1y_1 + \delta_1 + x_1\delta_1 + y_1\delta_1$$
  

$$+ x_1y_1\delta_1$$

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$$+ x_1y_1\delta_1$$
  

$$b = (1 + x_1 + y_1 + x_1y_1 + \delta_1 + x_1\delta_1 + y_1\delta_1$$
  

$$+ x_1y_1\delta_1)(10 + z_1)(1 + \delta_2)$$

b

 $= 10 + 10x_{1} + 10y_{1} + 10x_{1}y_{1}$  $+ 10\delta_{1} + 10x_{1}\delta_{1} + 10y_{1}\delta_{1} + 10x_{1}y_{1}\delta_{1}$  $+ z_{1} + x_{1}z_{1} + y_{1}z_{1} + x_{1}y_{1}z_{1}$  $+ \delta_{1}z_{1} + x_{1}\delta_{1}z_{1} + y_{1}\delta_{1}z_{1} + x_{1}y_{1}\delta_{1}z_{1}$  $+ 10\delta_{2} + 10x_{1}\delta_{2} + 10y_{1}\delta_{2} + 10x_{1}y_{1}\delta_{2}$  $+ 10\delta_{1}\delta_{2} + 10x_{1}\delta_{1}\delta_{2} + 10y_{1}\delta_{1}\delta_{2} + 10x_{1}y_{1}\delta_{1}\delta_{2}$  $+ z_{1}\delta_{2} + x_{1}z_{1}\delta_{2} + y_{1}z_{1}\delta_{2} + x_{1}y_{1}z_{1}\delta_{2}$  $+ \delta_{1}z_{1}\delta_{2} + x_{1}\delta_{1}z_{1}\delta_{2} + y_{1}\delta_{1}z_{1}\delta_{2} + x_{1}y_{1}\delta_{1}z_{1}\delta_{2}$ 

b

$$= 10 + 10x_{1} + 10y_{1} + 10x_{1}y_{1} + 10\delta_{1} + 10x_{1}\delta_{1} + 10y_{1}\delta_{1} + 10x_{1}y_{1}\delta_{1} + z_{1} + x_{1}z_{1} + y_{1}z_{1} + x_{1}y_{1}z_{1} + \delta_{1}z_{1} + x_{1}\delta_{1}z_{1} + y_{1}\delta_{1}z_{1} + x_{1}y_{1}\delta_{1}z_{1} + 10\delta_{2} + 10x_{1}\delta_{2} + 10y_{1}\delta_{2} + 10x_{1}y_{1}\delta_{2} + 10\delta_{1}\delta_{2} + 10x_{1}\delta_{1}\delta_{2} + 10y_{1}\delta_{1}\delta_{2} + 10x_{1}y_{1}\delta_{1}\delta_{2} + z_{1}\delta_{2} + x_{1}z_{1}\delta_{2} + y_{1}z_{1}\delta_{2} + x_{1}y_{1}z_{1}\delta_{2} + \delta_{1}z_{1}\delta_{2} + x_{1}\delta_{1}z_{1}\delta_{2} + y_{1}\delta_{1}z_{1}\delta_{2} + x_{1}y_{1}\delta_{1}z_{1}\delta_{2}$$

What are the bounds on the range and relative error of b?

b

$$= 10 + 10x_{1} + 10y_{1} + 10x_{1}y_{1} + 10\delta_{1} + 10x_{1}\delta_{1} + 10y_{1}\delta_{1} + 10x_{1}y_{1}\delta_{1} + z_{1} + x_{1}z_{1} + y_{1}z_{1} + x_{1}y_{1}z_{1} + \delta_{1}z_{1} + x_{1}\delta_{1}z_{1} + y_{1}\delta_{1}z_{1} + x_{1}y_{1}\delta_{1}z_{1} + 10\delta_{2} + 10x_{1}\delta_{2} + 10y_{1}\delta_{2} + 10x_{1}y_{1}\delta_{2} + 10\delta_{1}\delta_{2} + 10x_{1}\delta_{1}\delta_{2} + 10y_{1}\delta_{1}\delta_{2} + 10x_{1}y_{1}\delta_{1}\delta_{2} + z_{1}\delta_{2} + x_{1}z_{1}\delta_{2} + y_{1}z_{1}\delta_{2} + x_{1}y_{1}z_{1}\delta_{2} + \delta_{1}z_{1}\delta_{2} + x_{1}\delta_{1}z_{1}\delta_{2} + y_{1}\delta_{1}z_{1}\delta_{2} + x_{1}y_{1}\delta_{1}z_{1}\delta_{2}$$

What are the bounds on the range and relative error of b? This is computing  $x \times y \times z!!$ 

b

 $= 10 + 10x_{1} + 10y_{1} + 10x_{1}y_{1}$  $+ 10\delta_{1} + 10x_{1}\delta_{1} + 10y_{1}\delta_{1} + 10x_{1}y_{1}\delta_{1}$  $+ z_{1} + x_{1}z_{1} + y_{1}z_{1} + x_{1}y_{1}z_{1}$  $+ \delta_{1}z_{1} + x_{1}\delta_{1}z_{1} + y_{1}\delta_{1}z_{1} + x_{1}y_{1}\delta_{1}z_{1}$  $+ 10\delta_{2} + 10x_{1}\delta_{2} + 10y_{1}\delta_{2} + 10x_{1}y_{1}\delta_{2}$  $+ 10\delta_{1}\delta_{2} + 10x_{1}\delta_{1}\delta_{2} + 10y_{1}\delta_{1}\delta_{2} + 10x_{1}y_{1}\delta_{1}\delta_{2}$  $+ z_{1}\delta_{2} + x_{1}z_{1}\delta_{2} + y_{1}z_{1}\delta_{2} + x_{1}y_{1}z_{1}\delta_{2}$  $+ \delta_{1}z_{1}\delta_{2} + x_{1}\delta_{1}z_{1}\delta_{2} + y_{1}\delta_{1}z_{1}\delta_{2} + x_{1}y_{1}\delta_{1}z_{1}\delta_{2}$ 

 $b = 10 + 10x_1 + 10y_1 + 10x_1y_1 + 10x_1y_1 + 10\delta_1 + 10x_1\delta_1 + 10y_1\delta_1 + 10x_1y_1\delta_1 + z_1 + x_1z_1 + y_1z_1 + x_1y_1z_1 + \delta_1z_1 + x_1\delta_1z_1 + y_1\delta_1z_1 + x_1y_1\delta_1z_1 + 10\delta_2 + 10x_1\delta_2 + 10y_1\delta_2 + 10x_1y_1\delta_2 + 10\delta_1\delta_2 + 10x_1\delta_1\delta_2 + 10y_1\delta_1\delta_2 + 10x_1y_1\delta_1\delta_2 + z_1\delta_2 + x_1z_1\delta_2 + y_1z_1\delta_2 + x_1y_1z_1\delta_2$ 

 $+\delta_1 z_1 \delta_2 + x_1 \delta_1 z_1 \delta_2 + y_1 \delta_1 z_1 \delta_2 + x_1 y_1 \delta_1 z_1 \delta_2$ 

 $= 10 + 10x_{1} + 10y_{1} + 10x_{1}y_{1}$  $+ 10\delta_{1} + 10x_{1}\delta_{1} + 10y_{1}\delta_{1} + 10x_{1}y_{1}\delta_{1}$  $+ z_{1} + x_{1}z_{1} + y_{1}z_{1} + x_{1}y_{1}z_{1}$  $+ \delta_{1}z_{1} + x_{1}\delta_{1}z_{1} + y_{1}\delta_{1}z_{1} + x_{1}y_{1}\delta_{1}z_{1}$  $+ 10\delta_{2} + 10x_{1}\delta_{2} + 10y_{1}\delta_{2} + 10x_{1}y_{1}\delta_{2}$  $+ 10\delta_{1}\delta_{2} + 10x_{1}\delta_{1}\delta_{2} + 10y_{1}\delta_{1}\delta_{2} + 10x_{1}y_{1}\delta_{1}\delta_{2}$  $+ z_{1}\delta_{2} + x_{1}z_{1}\delta_{2} + y_{1}z_{1}\delta_{2} + x_{1}y_{1}z_{1}\delta_{2}$  $+ \delta_{1}z_{1}\delta_{2} + x_{1}\delta_{1}z_{1}\delta_{2} + y_{1}\delta_{1}z_{1}\delta_{2} + x_{1}y_{1}\delta_{1}z_{1}\delta_{2}$ 

Can contribute  $\pm 1.1920929 \times 10^{-10}$  to final range of b

h

Can contribute ±2 to final range of b h  $= 10 + 10x_1 + 10y_1 + 10x_1y_1$  $+10\delta_{1} + 10x_{1}\delta_{1} + 10y_{1}\delta_{1} + 10x_{1}y_{1}\delta_{1}$  $+ z_1 + x_1 z_1 + y_1 z_1 + x_1 y_1 z_1$  $+\delta_1 z_1 + x_1 \delta_1 z_1 + y_1 \delta_1 z_1 + x_1 y_1 \delta_1 z_1$  $+10\delta_{2}+10x_{1}\delta_{2}+10y_{1}\delta_{2}+10x_{1}y_{1}\delta_{2}$  $+10\delta_{1}\delta_{2} + 10x_{1}\delta_{1}\delta_{2} + 10y_{1}\delta_{1}\delta_{2} + 10x_{1}y_{1}\delta_{1}\delta_{2}$  $+ z_1 \delta_2 + x_1 z_1 \delta_2 + y_1 z_1 \delta_2 + x_1 y_1 z_1 \delta_2$  $+\delta_1 z_1 \delta_2 + x_1 \delta_1 z_1 \delta_2 + y_1 \delta_1 z_1 \delta_2 + x_1 y_1 \delta_1 z_1 \delta_2$ 

Can contribute  $\pm 1.1920929 \times 10^{-10}$  to final range of b

b

$$= 10 + 10x_{1} + 10y_{1} + 10x_{1}y_{1} + 10\delta_{1} + 10x_{1}\delta_{1} + 10y_{1}\delta_{1} + 10x_{1}y_{1}\delta_{1} + z_{1} + x_{1}z_{1} + y_{1}z_{1} + x_{1}y_{1}z_{1} + \delta_{1}z_{1} + x_{1}\delta_{1}z_{1} + y_{1}\delta_{1}z_{1} + x_{1}y_{1}\delta_{1}z_{1} + 10\delta_{2} + 10x_{1}\delta_{2} + 10y_{1}\delta_{2} + 10x_{1}y_{1}\delta_{2} + 10\delta_{1}\delta_{2} + 10x_{1}\delta_{1}\delta_{2} + 10y_{1}\delta_{1}\delta_{2} + 10x_{1}y_{1}\delta_{1}\delta_{2} + z_{1}\delta_{2} + x_{1}z_{1}\delta_{2} + y_{1}z_{1}\delta_{2} + x_{1}y_{1}z_{1}\delta_{2} + \delta_{1}z_{1}\delta_{2} + x_{1}\delta_{1}z_{1}\delta_{2} + y_{1}\delta_{1}z_{1}\delta_{2} + x_{1}y_{1}\delta_{1}z_{1}\delta_{2}$$

 $b = 10 + 10x_1 + 10y_1 + 10x_1y_1 + 10\delta_1 + z_1$  $+ x_1z_1 + y_1z_1 + x_1y_1z_1 + 10\delta_2$ 

h  $= 10 + 10x_1 + 10y_1 + 10x_1y_1$  $+10\delta_1 + 10x_1\delta_1 + 10y_1\delta_1 + 10x_1y_1\delta_1$  $+ z_1 + x_1 z_1 + y_1 z_1 + x_1 y_1 z_1$  $+\delta_1 z_1 + x_1 \delta_1 z_1 + y_1 \delta_1 z_1 + x_1 y_1 \delta_1 z_1$  $+ 10\delta_2 + 10x_1\delta_2 + 10y_1\delta_2 + 10x_1y_1\delta_2$  $10\delta_{1}\delta_{2} + 10\kappa_{1}\delta_{1}\delta_{2} + 10\gamma_{1}\delta_{1}\delta_{2} + 10\kappa_{1}\gamma_{1}\delta_{1}\delta_{2}$  $+ z_1 \delta_2 + x_1 z_1 \delta_2 + y_1 z_1 \delta_2 + x_1 y_1 z_1 \delta_2$ 

 $b = 10 + 10x_1 + 10y_1 + 10x_1y_1 + 10\delta_1 + z_1$  $+ x_1z_1 + y_1z_1 + x_1y_1z_1 + 10\delta_2 + \zeta_1 \qquad |\zeta_1| \le 0.0015$ 

# An added bonus

- Can use methods from approximation theory to make our technique applicable to algorithms including any elementary functions (e.g. Sine/Cosine/Sqrt)
  - These methods approximate an elementary function using a polynomial and an extra term bounding the error of the approximation

b

$$= 10 + 10x_{1} + 10y_{1} + 10x_{1}y_{1} + 10\delta_{1} + 10x_{1}\delta_{1} + 10y_{1}\delta_{1} + 10x_{1}y_{1}\delta_{1} + z_{1} + x_{1}z_{1} + y_{1}z_{1} + x_{1}y_{1}z_{1} + \delta_{1}z_{1} + x_{1}\delta_{1}z_{1} + y_{1}\delta_{1}z_{1} + x_{1}y_{1}\delta_{1}z_{1} + 10\delta_{2} + 10x_{1}\delta_{2} + 10y_{1}\delta_{2} + 10x_{1}y_{1}\delta_{2} + 10\delta_{1}\delta_{2} + 10x_{1}\delta_{1}\delta_{2} + 10y_{1}\delta_{1}\delta_{2} + 10x_{1}y_{1}\delta_{1}\delta_{2} + z_{1}\delta_{2} + x_{1}z_{1}\delta_{2} + y_{1}z_{1}\delta_{2} + x_{1}y_{1}z_{1}\delta_{2} + \delta_{1}z_{1}\delta_{2} + x_{1}\delta_{1}z_{1}\delta_{2} + y_{1}\delta_{1}z_{1}\delta_{2} + x_{1}y_{1}\delta_{1}z_{1}\delta_{2}$$

 $b = 10 + 10x_1 + 10y_1 + 10x_1y_1 + 10\delta_1 + z_1$  $+ x_1z_1 + y_1z_1 + x_1y_1z_1 + 10\delta_2 + \zeta_1 \qquad |\zeta_1| \le 0.0015$ 

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We get this for free!!

#### Tests

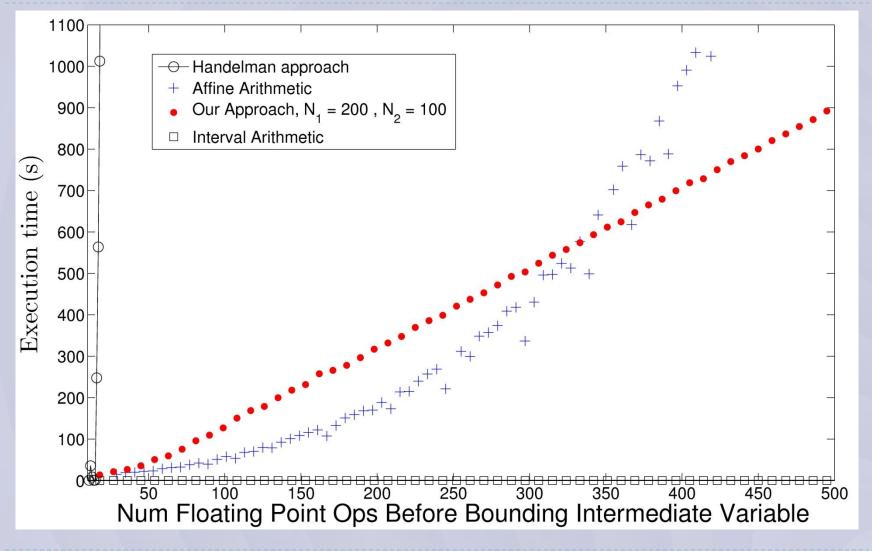
#### 5x5 Successive over relaxation

Real algorithm to find the solution to a system of linear equations of form Ax = b

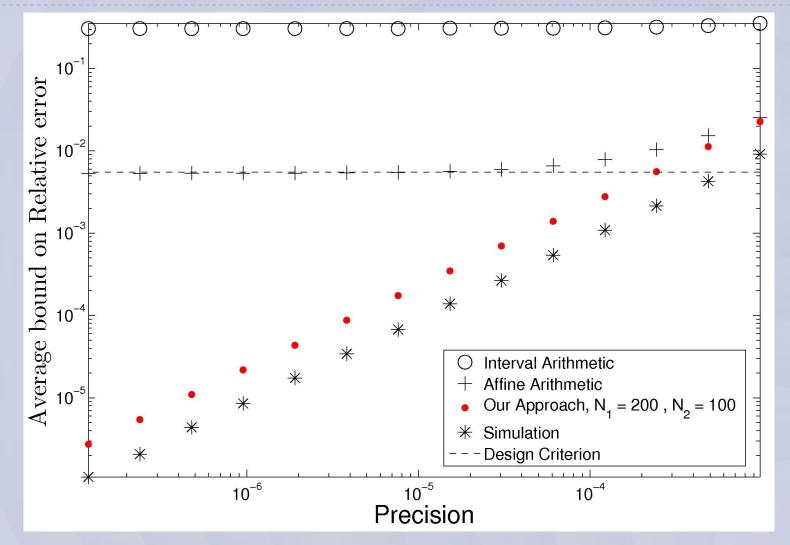
for 
$$k = 1; k \le 8; k + 4\mathbf{0}$$
  
for  $j = 1; j \le 5; j + 4\mathbf{0}$   
 $x^{j} = (1 - w)x^{j} + \frac{w}{A(j)^{j}}(b_{j} - \sum_{i=1, i \neq j}^{5} A(j)^{i}x^{i})$   
and for

end for

# Scalability: Execution time vs #operations



# Quality of bounds: Relative error vs precision



# Hardware use

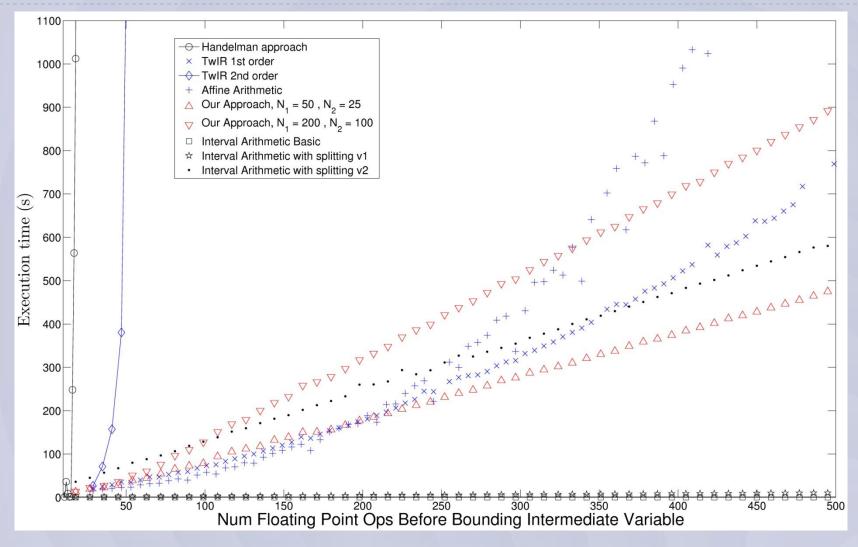
Method	Exponent	Mantissa	Slice	Slice	Frequency
	(# bits)	(# bits)	Regs	LUTs	(MHz)
Simulation	8	11	3562	3012	330
Our Approach	8	13	4261	3647	330
Affine Arithmetic	8	18	6606	5368	300
IA	$\infty$	$\infty$	$\infty$	$\infty$	N/A
IEEE Single Precision	8	24 – –	8407	6815	280
IEEE Double Precision	11	53	27200	22066	251

# Summary

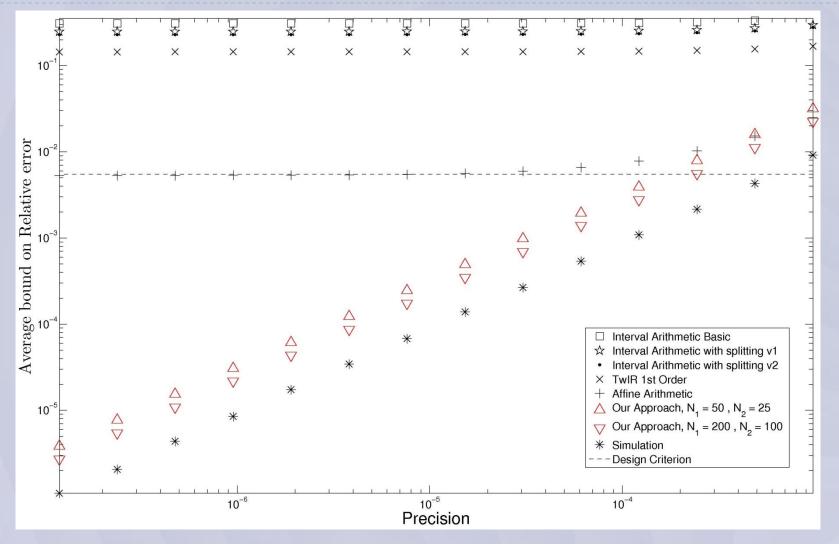
- Word-length optimisation can significantly improve hardware
- Need scalable analysis techniques to apply word-length optimisation on larger, more complex algorithms
- Our paper describes a simple set of algorithms to obtain tight bounds within a scalable execution time
  - Can use >80% fewer slice registers than IEEE double precision arithmetic
  - Can use >30% fewer slice registers than competing methods.
  - Can create hardware that is guaranteed to meet design criteria that is not possible using alternative methods

# Thank you for listening

# Scalability: Execution time vs #operations



# Quality of bounds: Relative error vs precision



# Quality of bounds vs execution time

